

# $\lambda$ Prolog: An Extended Logic Programming Language

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The logic programming language  $\lambda$ Prolog is an extension of conventional Prolog in several different directions. These extensions provide higher-order functions,  $\lambda$ -terms, a polymorphic typing discipline, modules, and a mechanism for providing secure abstract datatypes. Our original goal in developing  $\lambda$ Prolog was to understand the essential logical and proof theoretic nature of these extensions. This work has led us to describe a class of formulas called *higher-order hereditary Harrop formulas* which play a role in  $\lambda$ Prolog that is similar to the role of positive Horn clauses in Prolog. This extended class of formulas permits stronger forms of logical reasoning than can be found in Prolog. For example, it allows universal quantification and implications into the bodies of program clauses as well as some forms of higher-order quantification. Higher-order hereditary Harrop formulas, therefore, significantly extend positive Horn clauses, and, consequently,  $\lambda$ Prolog significantly enriches Prolog.

Part of our efforts have also been directed at understanding how this enrichment can be used to write programs. In this direction we have been investigating how the mechanisms of  $\lambda$ Prolog can be used to implement theorem provers, program transformers, and natural language understanding systems.

We have also implemented most features of higher-order hereditary Harrop formulas in a prototype interpreter, called  $\lambda$ Prolog V2.6. This interpreter comprises roughly 4100 lines of C-Prolog code and has been distributed to about 30 sites in North America and Europe since August 1987. The performance of this interpreter leaves much to be desired, owing, in part, to the underlying implementation language and, to a much larger extent, to the fact that efficiency has not been a major concern in this experimental version. Despite this drawback, the system has been used by us and others for serious experimentation and prototype implementations. Two ongoing Ph.D. theses make use of it as their primary implementation language.

Below is an outline of the major research aspects of this effort to date.

**Logic programming foundations** The theory of higher-order Horn clauses is outlined in [6] and given in detail in [9]. The higher-order unification process of [2] plays a central role in this analysis. The first-order theory of a language which includes implications in the bodies of program clauses is contained in [4]. Higher-order hereditary Harrop formulas were introduced in [8] to describe program clauses which are both higher-order and contain

implications and universal quantification in clause bodies.

**Modules for logic programming** The extended use of implications and universal quantifiers permits notions of modules and abstract datatypes to be supported. See [4] and [8].

**Theorem prover implementation**  $\lambda$ Prolog has proved to be a very natural implementation environment for writing tactic style theorem provers [1].

**Program transformation** Extending the work of [3], we have experimented with implementing several program transformation algorithms. The relationship between the actual implementation code in  $\lambda$ Prolog and the semantics of the programs being transformed seems very tight. We hope this will permit us to establish formal correctness proofs for these transformers [7].

**Computational linguistics** Natural language understanding systems must bring together syntactic and semantic processing. Very often syntactic processing can be identified with first-order operations while semantic processing is often higher-order.  $\lambda$ Prolog seems to offer a good environment for bringing these two kinds of processing together [5].

## References

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